

Determiners and generalized quantifiers

1 The

- A *the*-phrase denotes an referential individual of type e . The definite determiner *the* is of type $\langle et, e \rangle$: it combines with a common noun of type $\langle e, t \rangle$ to return an referential individual of type e .¹

- (2) a. $\llbracket \text{the cat} \rrbracket = \iota x_e [\text{cat}(x)]$
(The unique entity x such that x is a cat)
- b. $\llbracket \text{the} \rrbracket = \lambda P_{\langle e, t \rangle} . \iota x_e [P(x)]$
($\llbracket \text{the} \rrbracket$ applies to a predicate $P_{\langle e, t \rangle}$ and returns the unique entity x s.t. $P(x)$ is true.)
- c. $\llbracket \text{The cat snores} \rrbracket = \llbracket \text{snores} \rrbracket (\llbracket \text{the cat} \rrbracket)$
 $= (\lambda y_e . \text{snores}(y)) (\iota x_e [\text{cat}(x)])$
 $= \text{snores}(\iota x_e [\text{cat}(x)])$

The definite determiner *the* presupposes uniqueness.

- (3) a. [Pointing at one cat], the cat snores.
b. [Pointing at two cats], # the cat snores.

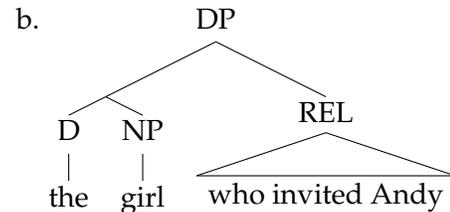
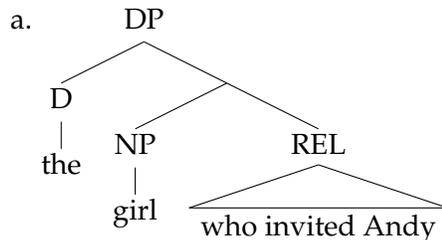
This presupposition is introduced by the ι -operator:

- (4) $\iota x_e [P(x)]$ is defined iff there exists exactly one x such that $P(x) = 1$.

- Adding relative clauses (REL)

Discussion: Which of the following (simplified) trees correctly describes the structure of the definite description “the girl who invited Andy”? [In other words, does the relative clause “who invited Andy” modify “girl” or “the girl”?] Why?

- (5) the girl who invited Andy

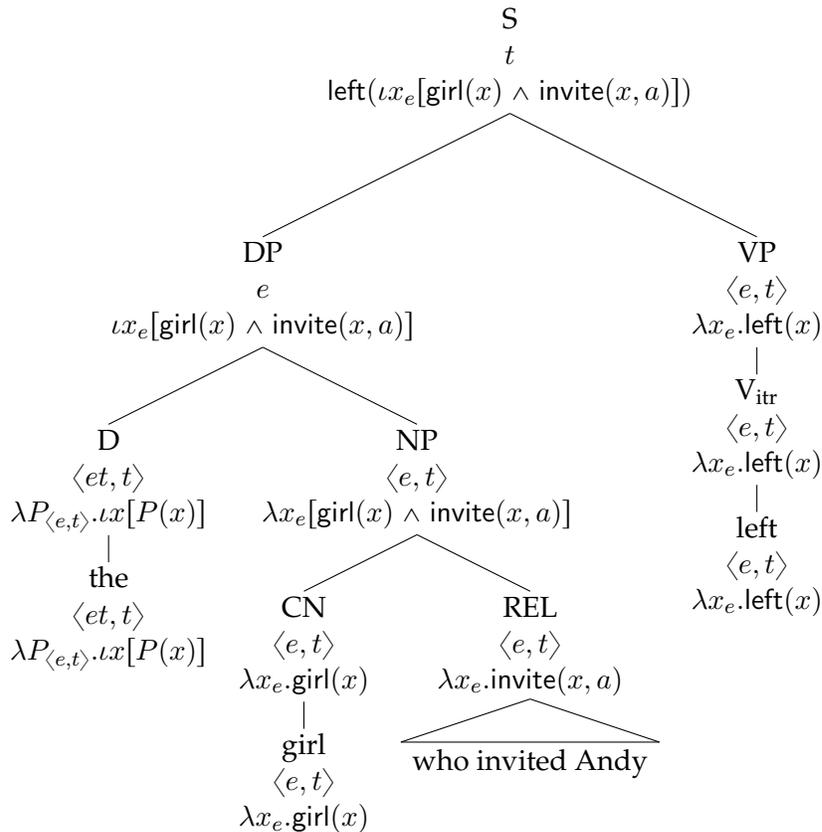


¹The introduced lexicon of *the* is highly simplified. It doesn't explain why we can say “the boys”, since it is nonsense to say “the unique boys”. A more accurate lexical entry of *the* is as follows:

- (1) $\llbracket \text{the} \rrbracket = \lambda P_{\langle e, t \rangle} . \exists x [P(x) \wedge \forall y [P(y) \rightarrow y \leq x]]$.
 $\iota x_e [P(x) \wedge \forall y [P(y) \rightarrow y \leq x]]$

- **Exercise:** Compose the following sentence:

(6) The girl who invited Andy left.



2 Generalized quantifiers and quantificational determiners

- Recall: How would you translate the following English sentences in set-theoretical notations?

- (7) a. Some cat meows.
 b. Every cat meows.
 c. No cat meows.

- In set-theoretical notations, a quantificational determiners like *some* and *every* is defined as a relation between two sets of entities. Then, how should we define phrases like *some/every/no cat* and determiners like *some/every/no* in λ -notations?

2.1 Generalized quantifiers

- Quantificational DPs (e.g. *everything*, *something*, *nothing*, *every cat*, *some cat*, *no cat*) are not individuals:

- (8) *Law of Contradiction*
- a. Mary is coming and Mary is not coming. (Contradiction)
 b. Someone is coming and someone is not coming. (Not contradiction)

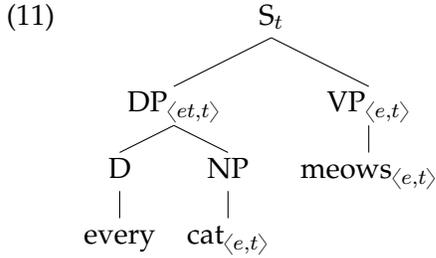
(9) *Law of Excluded middle*

- a. Mary is coming or Mary is not coming. (Tautology)
- b. Every is coming or everyone is not coming. (Not tautology)

(10) *Only an e-type NP can normally license a singular discourse pronoun.*

- a. John /the man/ a man walked in. He looked tired.
- b. Every man /no man/ more than one man walked in. *He looked tired.

- We treat quantificational DPs as second-order functions of type $\langle et, t \rangle$, called **generalized quantifiers**. In (11), *meows* is an argument of *every cat*.



- (12) a. $\llbracket \text{every cat} \rrbracket = \lambda P_{\langle e, t \rangle}. \forall x [\text{cat}(x) \rightarrow P(x)]$
 b. $\llbracket \text{every cat meows} \rrbracket = \llbracket \text{every cat} \rrbracket (\llbracket \text{meows} \rrbracket)$
 $= (\lambda P_{\langle e, t \rangle}. \forall x [\text{cat}(x) \rightarrow P(x)]) (\lambda y_e. \text{meows}(y))$
 $= \forall x [\text{cat}(x) \rightarrow \text{meows}(x)]$

- (13) a. $\llbracket \text{some cat} \rrbracket = \lambda P_{\langle e, t \rangle}. \exists x [\text{cat}(x) \wedge P(x)]$
 b. $\llbracket \text{some cat meows} \rrbracket = \exists x [\text{cat}(x) \wedge \text{meows}(x)]$

- (14) a. $\llbracket \text{no cat} \rrbracket = \lambda P_{\langle e, t \rangle}. \neg \exists x [\text{cat}(x) \wedge P(x)]$
 b. $\llbracket \text{no cat meows} \rrbracket = \neg \exists x [\text{cat}(x) \wedge \text{meows}(x)]$

- Individuals (of type e) can also be shifted into generalized quantifiers via *type-lifting*.

- (15) a. $\llbracket \text{Kitty} \rrbracket = j$
 b. $\text{LIFT}(\llbracket \text{Kitty} \rrbracket) = \lambda P_{\langle e, t \rangle}. P(j)$
 c. $(\text{LIFT}(\llbracket \text{Kitty} \rrbracket))(\llbracket \text{meows} \rrbracket) = (\lambda P_{\langle e, t \rangle}. P(j))(\lambda x. \text{meows}(x))$
 $= (\lambda x. \text{meows}(x))(j)$
 $= \text{meows}(j)$

2.2 Quantificational determiners

- The determiner *every* combines with a common noun of type $\langle e, t \rangle$ to return a generalized quantifier of type $\langle et, t \rangle$. Therefore, its type is quite complex: $\langle et, \langle et, t \rangle \rangle$.

- (16) a. $\llbracket \text{every} \rrbracket = \lambda Q_{\langle e, t \rangle}. \lambda P_{\langle e, t \rangle}. \forall x [Q(x) \rightarrow P(x)]$
 b. $\llbracket \text{some} \rrbracket = \lambda Q_{\langle e, t \rangle}. \lambda P_{\langle e, t \rangle}. \exists x [Q(x) \wedge P(x)]$
 c. $\llbracket \text{no} \rrbracket = \lambda Q_{\langle e, t \rangle}. \lambda P_{\langle e, t \rangle}. \neg \exists x [Q(x) \wedge P(x)]$